Introduction to

Applied Topology

Lecture Notes

Introduction to Applied Topology In the next 5 weeks, we will (roughly) cover: - invariants: culer characteristic, (co) homology, persistence diagrams - notions of stability: classical i recent thms i proofs functionals of persistence diagrams manifold learning -applications: topology of random spaces multiparameter persistence sheaf theory - generalisations: other applications ooo What we will not cover : homotopy limit theorems Note that these note will be expanded during the course

Preliminacies

Topological space - very general - defined in terms of neighborhoods -lead to many pathological examples Other types of spaces : Manifolds, stratified spaces, metric spaces, ... 1st model : Simplicial complex Def: A k-simplex is the convex combination of (k+1) - points • 0 Point O-simplex edge triangle tetrchedron 3-simplex 2-simplex 1-simple

If points are embedded in Euclidean space then a simplex is just the convex hall

Mote: Simplices can be abstract, met recessarily embedded (we will see this later)

One simplex is not terribly interesting - consider simplicial complexes. Def: A simplicial complex A is a set of simplicies Eog such that i) if del i red then red 2) if an az + of then an az es Examples . not simplicial complexes simplicial complex We will introduce related concepts as we need them, eg. carrier, closure, star, link,... Note: A simplicial complex is often the easiest to work with in a computer, but for hand computed examples we will often use cellular complexes

Def: A cellular complex consists of k-cells i attaching maps (how to glue the cell to lower dimensional cells Example: 2 cells are discs. cellulat Simplicial There are other models: &-complexes, simplicial sets Here we will mainly use simplicial complexes (with celluler for some small examples) Note: Simplicial and cellular Simplicial is always cellular Ж cellular is "equivalent" to simplicial (we will see exactly how later but one can so belivide a cells into simplices.

Equivalences

We will see many type of equivalances however, topological invariants are invariant under continuous transformations (this could be any continuous functions, homeomorphisms, diffeomorphisms, etc)

For our purposes, if f is a homeomorphism

for f: X -> Y then we consider

X~Y equivalent. As we introduce these concepts we will place them into

context (stronger vs. weaker)

Euler Characteristic

One of the most lasic i most ubiquitous topological invariants.

For surfaces (Euler)

 $\chi = |\chi| - |\Xi| + |E|$



<u>General</u> formula $\chi(k) = \sum_{i=0}^{d} (-i)^{i}(\# k-simplices in D)$

Facts: - Topological invariant - Easy to compute - Ubiquitous (topology, differential geometry, category theory, random fields, ...)

Chain Complexes

We are after computable invariants. Much of what we will use will be based on linear algebra. Our starting point is the chain complex.

We will first describe the most relevant example before formally defining it.



Define Col(D) := IIKal

The d-th chain group Cd(D) is the vector space where each simplex is its own dimension



Note: In everything we consider, the chain group is really the chain vector space.

* Why is it called the chain group?

We can add simplices together (or more precisely, we can take linear combinations of simplices)

Del A k-chain is the linear combination of K-simplicies



or in other words, it is a vector indexed by the k-simplicies.

So gar, we have only defined chains

what is a chain complex?

Recall : Simplicies (or cells) of different dimensions are related (if I is a face of a simplex tink then I must be in K) This gives rise to maps dr. Cx -> Ck-1 Hk dx is called the k-th boundary operator. It describes how a k-simplex maps to a (K-1) - chain. Example (over Z2) Ja 2(a) -> a+b • C. 2(abc) = abtbctac

In general, we will have orientation (-1.) as well

Notice that dx takes a k-simplex i returns a (k-1) chain. Since it is linear, more generally it takes a k-chain ; returns a (k-1)-chain.



Note: vertices map to 0, 2(1) = 0







Since de: Ce-> Ce-1, the chain complex is •••• $C_{k+1} \xrightarrow{\partial k} C_k \xrightarrow{\partial k} C_{k-1} \xrightarrow{\partial k-1} \cdots \xrightarrow{\partial k} C_1 \xrightarrow{\partial 1} C_0 \xrightarrow{\partial$ The collection E_{c_k} , $\partial_c \mathcal{J}_{ce2}$ is the chain complex. Property: A chain complex must also have the property $\partial_{k} \circ \partial_{k+1} = 0$ a this is the) this is the key condition Alternatively im OKHI E ker OK Example: a b. C O(abc) = ab + ac + bc 2.2(abc)= 2(abrac+bc) $= \partial(ab) + \partial(ac) + \partial(bc)$ = a + b + a + c + b + c= 0 (arc=0, b+b=0, crc=0) Since we are over Zz

Def: An element of kerdix is called a k-cycle.



Exercise: Verify that 2.2 (D) =0, that the boundary of an empty tetrahedron (4 triangles) is O,

Def: The space of all k-cycles is called the cycle space.

Def: An element of im DK+1 ECK is called a k-boundary (because it bounds a k-cycle)

Obs: These are all vector spaces (though this holds only in our case)

Homology Given a chain complex ECK, Dr3 let Zu denote the space of cycles (kerde) BKHI denote the space of boundaries (in deti) The k-the homology group is $H_{\mathcal{K}}(S) = \frac{\ker \partial \kappa}{\operatorname{im} \partial \kappa_{*}}$ "All cycles which are not bounded" (In our case again a vector space) Def: The k-th Betti number is the rank of Hr Bre = rk(Hr) = rk (ker Dr) = = rk kerde - rkimdeti

Why is this well-clefined? Recall: drodki = 0, so in Deti Sker De SCE 1 1 1 1 Loudories cycles chains rk(inder) ≤ rk(kerdx) ≤ rk(Cr) Picture: Cr Zr Br Zĸ 2-(BK-1) How can we compute this? Gaussian elimination





•



Euler Characteristic ? Homology Recall X = Z (-i) i [Ka] Now: |Kdl=rk(ck) (since ck is just the identity matrix indexed by k-simplices) By the rank nullity theorem rk(cc) = rk(kerde) + rk(imde) So $\chi = \sum_{i=0}^{d} (-i)^{i} r k (ck) = \sum_{i=0}^{d} (-i)^{i} (r k (ker \partial_{k}) + r k (im \partial_{k}))$ = Z(-N' (-k (ker dx) - - k (im dx, N + rk(in d)) $= \sum_{i=0}^{d} (-i)^{i} r (\psi_{i}) = \sum_{i=0}^{d} (-i)^{i} B_{i}$ Oddly, computing alternating sums of Betti numbers is much easier than computing Betti numbers.

Persistent Homology

Until now, we have had only one space A

Now we consider a filtration

 $\Delta_1 \subseteq \Delta_2 \subseteq \ldots \subseteq \Delta_n$

What is the homology of a filtration?

Digression: What are some typical filtrations we will look at?

i) Functions on simplicial complexes

f: A-> R

(we assume for simplicity, that I is constant on each simplex) The fitration f="(-∞, x] is called a Lower-star fitration.

We require that for (-or, x] is a simplicial complex (Edelsbrunner & Harer call this a monotone function

We will not prove it here - but there is a closely related notion - the sublevel set filtration.

Intuition: Track homological features over the filtration. Example: 1 on vew comp. new comp 1 2 comp. verge. 2 comp. merge * Fratures can be born (rk(Hz)+1) Features can die (rk (HE)-1) Same example for H, (with applogies)

Note: There is more information than just +1/-1 of the rank Notice +/ ranks of components is the same but what happens if we track how long live! Key Fact: When 2 components/cycles merge, we kill the youngest one, ie the one which was born last. This is called the Elder rule (There is a good algebraic reason for this)



Formal Definitions

Given a filtration of simplicial complexes

 $\Delta_1 \subseteq \Delta_2 \subseteq \Delta_3 \subseteq \cdots \subseteq \Delta_n$

This gives rise to an increasing sequence of chain groups for all k

 $C_{\kappa}(\Delta_{1}) \subset C_{\kappa}(\Delta_{2}) \subset C_{\kappa}(\Delta_{3}) \subset \cdots \subset C_{\kappa}(\Delta_{n})$









Each colum is a chain complex & Cs denotes monomorphisms

Applying homology in each column $C_{\kappa}(\Delta_{1}) \subset C_{\kappa}(\Delta_{2}) \subset C_{\kappa}(\Delta_{3}) \subset \cdots \subset C_{\kappa}(\Delta_{m})$ $\mathcal{H}_{\mathcal{L}}(\mathcal{D}_{1}) \longrightarrow \mathcal{H}_{\mathcal{L}}(\mathcal{D}_{2}) \longrightarrow \mathcal{H}_{\mathcal{L}}(\mathcal{D}_{3}) \longrightarrow \mathcal{H}_{\mathcal{L}}(\mathcal{D}_{3})$ Note: Usually in topology we compute homology over & (ie we treat the entries in de as integers) but we will treat them as Z2 (or Ep for p prime)e This means HiclA:) are vector spaces ; Hr (Di) -> Hr (Di) are linear maps. Remark I: The linear maps are induced from the inclusions on the chain groups Ex: Check that the maps are well defined Hint: CK(Si) C> Cc(Si) $\Rightarrow Z_{\varepsilon}(D_{i}) \hookrightarrow Z_{\varepsilon}(D_{i})$ => Br (Di) C> Br (Dj)

Leuark 2: The linear maps are just matrices again Remark 3: Though the chain maps are monomorphisms, this is not the case for the linear maps. Def. A persistence module is the collection of vector spaces and imaps between them: EHE (Di) Siez & HE (Di) ->HE (Di) W; j This is just: $\mathcal{H}_{\mathcal{K}}(\Delta_{1}) \longrightarrow \mathcal{H}_{\mathcal{K}}(\Delta_{2}) \longrightarrow \mathcal{H}_{\mathcal{K}}(\Delta_{3}) \longrightarrow \mathcal{H}_{\mathcal{K}}(\Delta_{3})$ One thing which makes persistence useful Module en Diagran Barcode

Def: A barcode is a decomposition of a persistence module PLD $P_k(\Delta) \cong I_i^k(A) \cong I_2^k(A) \cong I_N^k(A)$ such that each bjetedj bj,djetk otherwise I;(t) = { o (These an called interval modules) $\frac{1}{1} r \left(H_{E}(\Delta_{S}) - 3 H_{E}(\Delta_{t}) \right) = \sum_{\delta=0}^{N} min I_{\delta}(t) \cap [S, t]$ So we decompose into "intervals" such that the rank of any map is the sum of intervals which "span" the map lie the interval contains the end points of the map) Why does this exist? Short answer: Gabriel's Thun (Representation theory) A quiver of the form admits a decomposition.

More constructive viewpoint:

Say we build an interval decomposition incrementally:

 $H_{\mathcal{K}}(\Delta_{1}) \xrightarrow{g_{12}} H_{\mathcal{K}}(\Delta_{2}) \xrightarrow{g_{23}} H_{\mathcal{K}}(\Delta_{3})$

Note: f.2 = f23 0 f2





Observation: All the information is contained in ranks of all the vector spaces and all maps. This makes sense since the rank determines (up to isomorphism) a vector spaces Computation Good news: No harder than ordinary homology Idea: Ordered Gaussian elimination Example: 0 2/8 406 * We will refer to simplices by function value > increasing 22 3567cogcle I are the intervals \odot l [1,3], [2,5], [4,6]6 OOT L [0,00] - un matched rows are infinite bars Carlsson-Zomorodian showed just include rows that we can restrict dx to which oreate cycles 9, 8 · 7 [] ---> [7,8]

This restriction makes computation easier but it is important to keep in mind aling this works. Algebraic interpretation: Treat coefficients as monomials in I voriable, w/ powers encoding time 0 3 0 5 ··· 7 8 -2 ··· 6 The power is the différence between edge time 's vertex "time" * Multiplication by t noves things forward in time Example: a b flat=1 flbt=2 S(as) = 3 $\partial(ab) = t^2 a + t b$ Notice: these polynomicle keep toock of "time" Caclsson-tomordian: Persistent hours logey is homology over nonomials w! field coefficients (Zp [+])

Upshot: Zp[t] is a principle ideal donain, so à persistence module is a module over a p.id & admits a decomposition into a free part linfinite bars) à torsion (finite bars) * these are precisely the interval modules. From baccodes to diagrams It will often be useful to present baccodes as a diagram. For each interval draw a point n/ the start point on the x-axis of the end point on the y-aris $[1,3], [2,5], [4,6] \longrightarrow$

Diagrans will be useful for stability.

But first,

I more interpretation:



Sitteration of length 3, this is represented by

e----e

8

Hence, we need 6 integers (ranks) This is called the can't function.

 $R: I \rightarrow Z$

R is an integer valued function on the space of all possible intervals



